Genome Assembly & Alignment Primer Michael Schatz

Sept 27, 2012 Beyond the Genome



@mike_schatz / #BTG2012



Outline

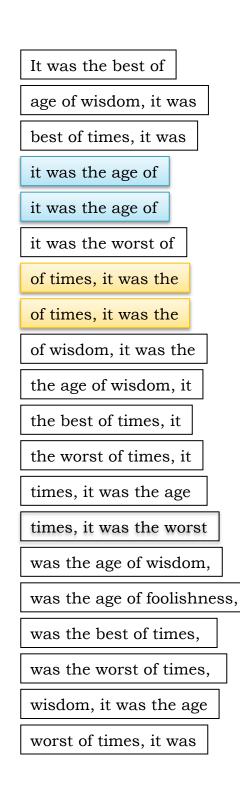
- I. Assembly by Analogy
- 2. Genome Assembly
 - I. Coverage, read length, repeats, and errors
 - 2. Genome assemblers & Assemblathon
- 3. Whole genome alignment

Shredded Book Reconstruction

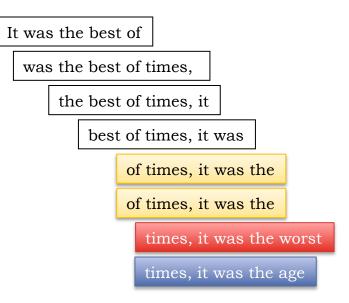
Dickens accidentally shreds the first printing of <u>A Tale of Two Cities</u>
 – Text printed on 5 long spools

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- How can he reconstruct the text?
 - 5 copies x 138, 656 words / 5 words per fragment = 138k fragments
 - The short fragments from every copy are mixed together
 - Some fragments are identical



Greedy Reconstruction



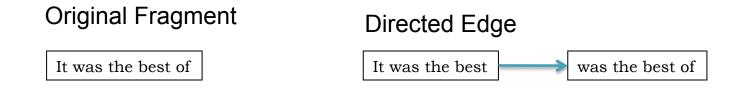
The repeated sequence make the correct reconstruction ambiguous

• It was the best of times, it was the [worst/age]

Model the assembly problem as a graph problem

de Bruijn Graph Construction

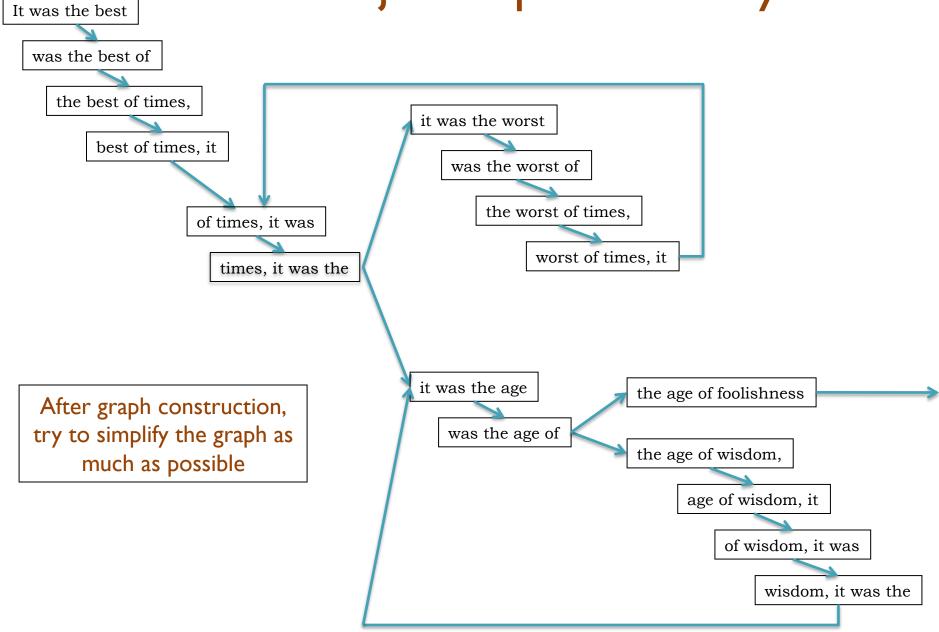
- $D_k = (V, E)$
 - V = All length-k subfragments (k < l)
 - E = Directed edges between consecutive subfragments
 - Nodes overlap by k-1 words



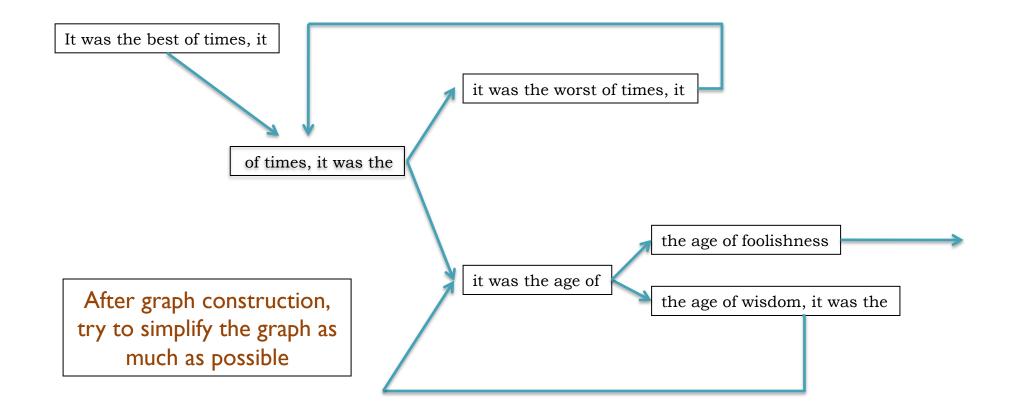
- Locally constructed graph reveals the global sequence structure
 - Overlaps between sequences implicitly computed

de Bruijn, 1946 Idury and Waterman, 1995 Pevzner, Tang, Waterman, 2001

de Bruijn Graph Assembly



de Bruijn Graph Assembly



Assembly Applications

Novel genomes





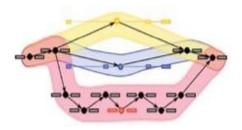
• Metagenomes





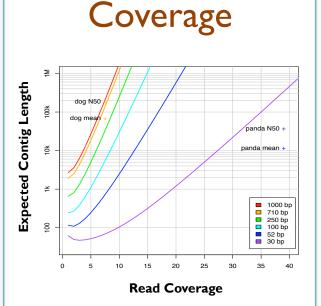
- Sequencing assays
 - Structural variations
 - Transcript assembly





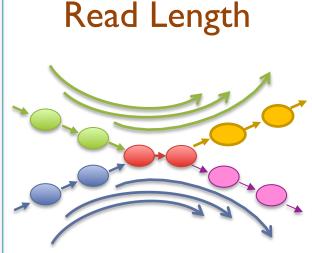
Like Dickens, we must computationally reconstruct a genome from short fragments

Ingredients for a good assembly



High coverage is required

- Oversample the genome to ensure every base is sequenced with long overlaps between reads
- Biased coverage will also fragment assembly





Reads & mates must be longer than the repeats

- Short reads will have *false overlaps* forming hairball assembly graphs
- With long enough reads, assemble entire chromosomes into contigs

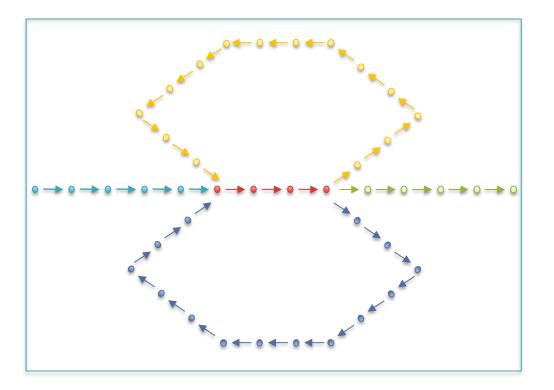
Errors obscure overlaps

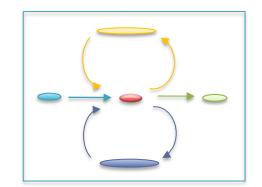
- Reads are assembled by finding kmers shared in pair of reads
- High error rate requires very short seeds, increasing complexity and forming assembly hairballs

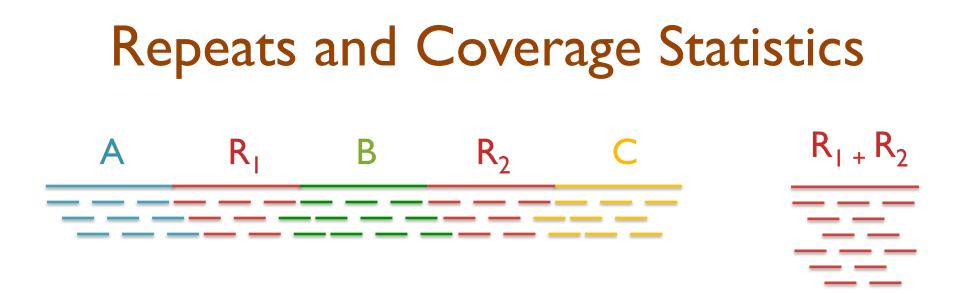
Current challenges in de novo plant genome sequencing and assembly Schatz MC, Witkowski, McCombie, WR (2012) *Genome Biology*. 12:243

Initial Contigs

- After constructing assembly graph, compress graph down to its non-branching initial contigs
 - Aka "unitigs", "unipaths"





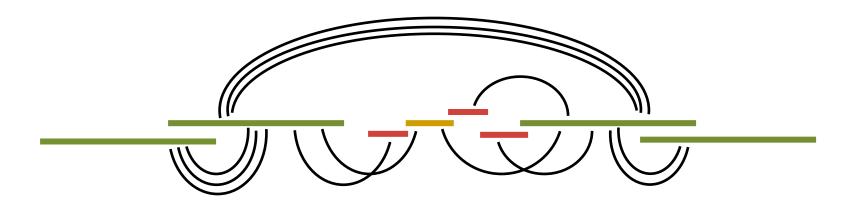


- If *n* reads are a uniform random sample of the genome of length *G*, we expect $k=n\Delta/G$ reads to start in a region of length Δ .
 - If we see many more reads than k (if the arrival rate is > A), it is likely to be a collapsed repeat
 - Requires an accurate genome size estimate

$$\Pr(X - copy) = \binom{n}{k} \left(\frac{X\Delta}{G}\right)^{k} \left(\frac{G - X\Delta}{G}\right)^{n-k} \qquad A(\Delta, k) = \ln\left(\frac{\Pr(1 - copy)}{\Pr(2 - copy)}\right) = \ln\left(\frac{\frac{(\Delta n/G)^{k}}{k!}e^{\frac{-\Delta n}{G}}}{\frac{(2\Delta n/G)^{k}}{k!}e^{\frac{-2\Delta n}{G}}}\right) = \frac{n\Delta}{G} - k\ln 2$$

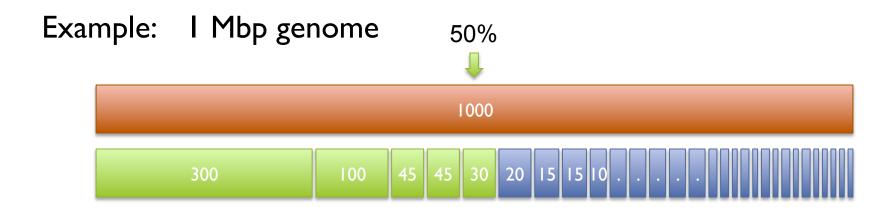
Scaffolding

- Initial contigs (aka unipaths, unitigs) terminate at
 - Coverage gaps: especially extreme GC regions
 - Conflicts: sequencing errors, repeat boundaries
- Iteratively resolve longest, 'most unique' contigs
 - Both overlap graph and de Bruijn assemblers initially collapse repeats into single copies
 - Uniqueness measured by a statistical test on coverage



N50 size

Def: 50% of the genome is in contigs larger than N50



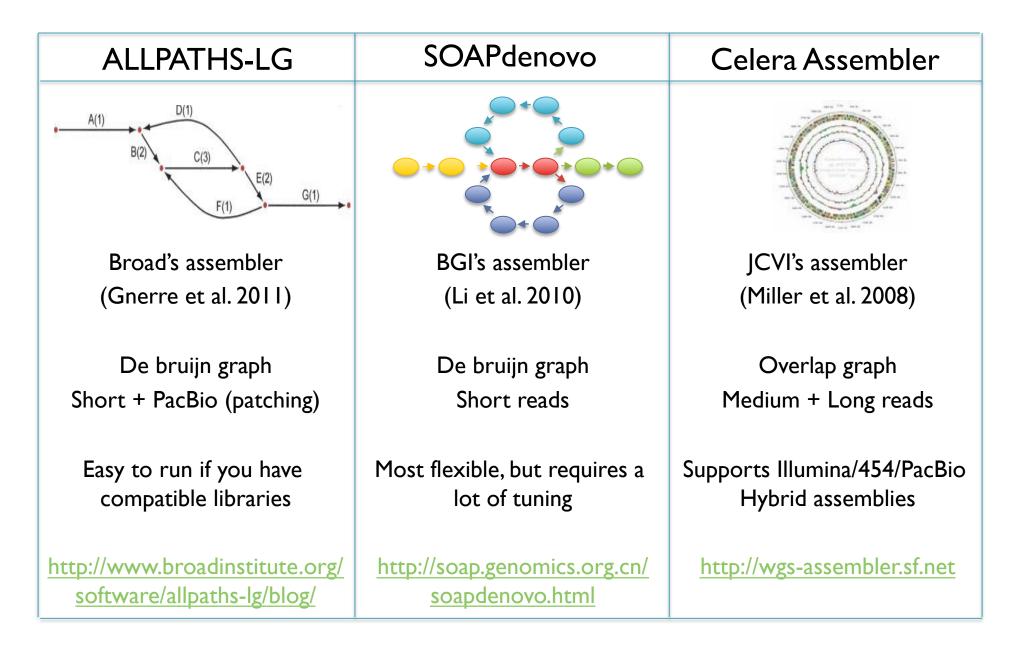
```
N50 size = 30 \text{ kbp}
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(300k+100k+45k+45k+30k = 520k \ge 500kbp)
```

Note:

N50 values are only meaningful to compare when base genome size is the same in all cases

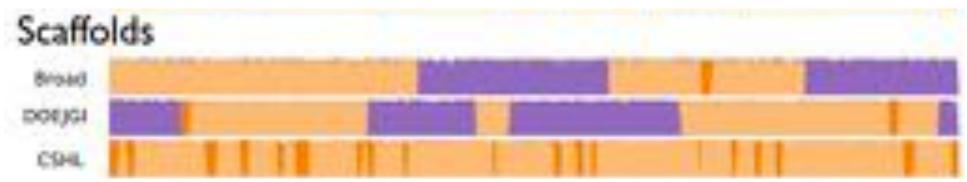
Assembly Algorithms

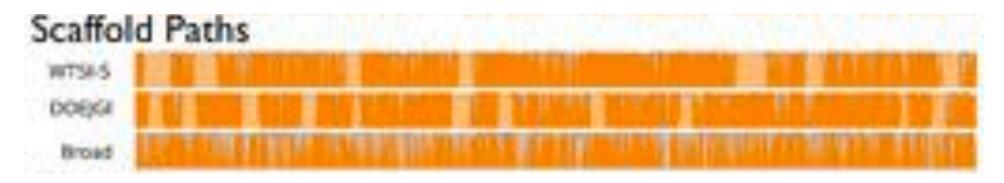


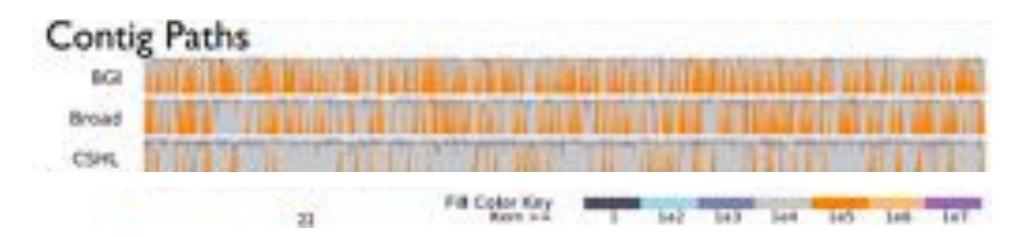


- Attempt to answer the question:
 "What makes a good assembly?"
- Organizers provided simulated sequence data
 - Simulated 100 base pair Illumina reads from simulated diploid organism
- 41 submissions from 17 groups
- Results demonstrate trade-offs assemblers must make

Assembly Results





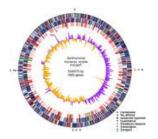


Final Rankings



- SOAPdenovo and ALLPATHS came out neck-and-neck followed closely behind by SGA, Celera Assembler, ABySS
- My recommendation for "typical" short read assembly is to use ALLPATHS

Assembly Summary



Assembly quality depends on

- I. Coverage: low coverage is mathematically hopeless
- 2. Repeat composition: high repeat content is challenging
- 3. Read length: longer reads help resolve repeats
- 4. Error rate: errors reduce coverage, obscure true overlaps
- Assembly is a hierarchical, starting from individual reads, build high confidence contigs/unitigs, incorporate the mates to build scaffolds
 - Extensive error correction is the key to getting the best assembly possible from a given data set
- Watch out for collapsed repeats & other misassemblies
 - Globally/Locally reassemble data from scratch with better parameters & stitch the 2 assemblies together

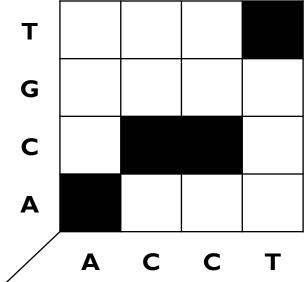


Whole Genome Alignment with MUMmer

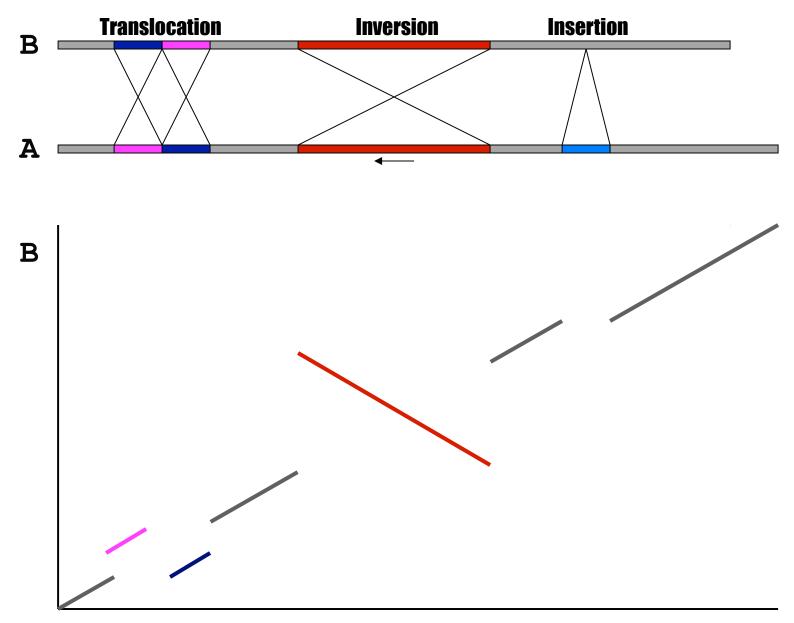
Slides Courtesy of Adam M. Phillippy <u>amp@umics.umd.edu</u>

WGA visualization

- How can we visualize *whole* genome alignments?
- With an alignment dot plot T $-N \times M$ matrix G• Let i = position in genome A• Let j = position in genome B• Fill cell (*i*,*j*) if A_i shows similarity to B_j A



 A perfect alignment between A and B would completely fill the positive diagonal





Seed and Extend

How can quickly find large alignments?

- I. Find short exact matches
 - using a suffix tree
- 2. Cluster exact matches
 - using size, gap and distance parameters
- 3. Extend clusters & report alignments
 - using modified Smith-Waterman algorithm

WGA example with **nucmer**

Yersina pestis CO92 vs. Yersina pestis KIM

- High nucleotide similarity, 99.86%
 - Two strains of the same species
- Extensive genome shuffling and highly repetitive
 - Global alignment will not work

nucmer -maxmatch CO92.fasta KIM.fasta
-maxmatch Find maximal exact matches (MEMs)

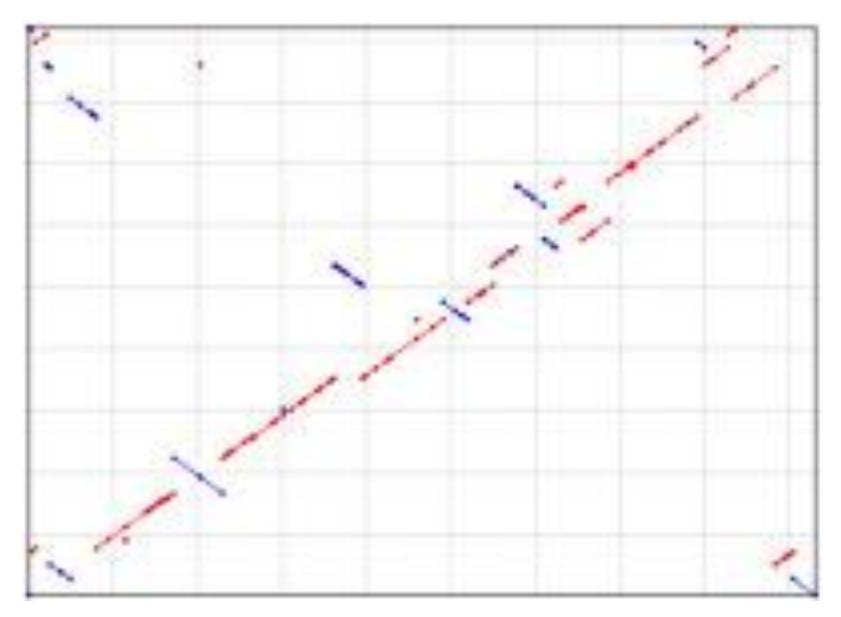
delta-filter -1 out.delta > out.filter.m
-n Many-to-many mapping

show-coords -r out.delta.m > out.coords

-r Sort alignments by reference position

mummerplot --large --layout out.delta.m

--large Large plot
--layout Nice layout for multi-fasta files



http://mummer.sourceforge.net

Thank You

http://schatzlab.cshl.edu @mike_schatz / #BTG2012

